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A Metaheuristic Optimization Algorithm for Binary Quadratic Problems

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$$\min_{x \in X} \sum_{i=1}^n \sum_{j=1}^n \sum_{k=1}^n \sum_{l=1}^n f_{ik} d_{jl} x_{ij} x_{kl} + \sum_{i=1}^n \sum_{j=1}^n c_{ij} x_{ij}$$

$$X = \{x \mid \sum_{j=1}^n x_{ij} = 1 \quad i \in N$$

$$\sum_{i=1}^n x_{ij} = 1 \quad j \in N$$

$$x_{ij} \in \{0, 1\} \quad i, j \in M\}$$



$$\sum_{i=1}^n \sum_{j=1}^n \sum_{k=1}^n \sum_{l=1}^n f_{ik} d_{jl} x_{ij} x_{kl} = \text{trace}(\mathbf{D}\mathbf{X}\mathbf{F}\mathbf{X}^T)$$



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$$\mathbf{F} = \mathbf{q}\mathbf{q}^T$$



$$\sum_{i=1}^n \sum_{j=1}^n \sum_{k=1}^n \sum_{l=1}^n f_{ik} d_{jl} x_{ij} x_{kl} = \text{trace}(\mathbf{DXFX}^T)$$

$$\mathbf{F} = \mathbf{q}\mathbf{q}^T$$

$$= \text{trace}(\mathbf{DXq}\mathbf{q}^T\mathbf{X}^T)$$

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$$= \text{trace}(\mathbf{y}^T\mathbf{Dy}) = \mathbf{y}^T\mathbf{Dy}$$



$$\min_{x \in X, y \in \mathbb{R}^n} y^T D y$$

subject to

$$y_i = \sum_{j=1}^n x_{ij} q_j \quad \forall i$$

$$\sum_{i=1}^n y_i = \sum_{j=1}^n q_j$$



$$\min_{x \in X, y, z \in \mathbb{R}^n} \mathbf{y}^T (\mathbf{D} + \text{Diag}(\mathbf{u})) \mathbf{y} - \mathbf{u}^T \mathbf{z}$$

subject to

$$y_i = \sum_{j=1}^n x_{ij} q_j \quad \forall i$$

$$z_i = \sum_{j=1}^n x_{ij} q_j^2 \quad \forall i$$

$$\sum_{i=1}^n y_i = \sum_{j=1}^n q_j$$



$$\min \mathbf{x}^T (\mathbf{D} + \text{diag}(\mathbf{u})) \mathbf{x} - \mathbf{u}^T \mathbf{x}$$

subject to

$$\sum_{i=1}^n x_i = k$$

iteration constraint

$$x_{iter} = 1$$



$$x = \begin{pmatrix} 1 \\ 0 \\ 1 \\ 1 \\ 0 \\ 0 \\ 0 \\ 1 \\ 0 \\ 0 \\ 0 \\ 1 \\ 1 \\ 1 \\ 0 \end{pmatrix}$$



$$x = \begin{pmatrix} 1 \\ 0 \\ 1 \\ 1 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 1 \\ 0 \end{pmatrix} \quad x_1 = \begin{pmatrix} 1 \\ 0 \\ 1 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 1 \\ 0 \end{pmatrix}$$



$$\mathbf{x} = \begin{pmatrix} 1 \\ 0 \\ 1 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 0 \\ 0 \\ 0 \\ 1 \\ 1 \\ 0 \end{pmatrix} \quad \mathbf{x}_1 = \begin{pmatrix} 1 \\ 0 \\ 1 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 0 \\ 0 \\ 0 \\ 1 \\ 1 \\ 0 \end{pmatrix} \quad \mathbf{x}_2 = \begin{pmatrix} 1 \\ 0 \\ 1 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 0 \\ 0 \\ 0 \\ 1 \\ 1 \\ 0 \end{pmatrix}$$



$$\mathbf{x} = \begin{pmatrix} 1 \\ 0 \\ 1 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 0 \\ 0 \\ 0 \\ 1 \\ 1 \\ 0 \end{pmatrix} \quad \mathbf{x}_1 = \begin{pmatrix} 1 \\ 0 \\ 1 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 0 \\ 0 \\ 0 \\ 1 \\ 1 \\ 0 \end{pmatrix} \quad \mathbf{x}_2 = \begin{pmatrix} 1 \\ 0 \\ 1 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 0 \\ 0 \\ 0 \\ 1 \\ 1 \\ 0 \end{pmatrix} \quad \mathbf{x}_3 = \begin{pmatrix} 1 \\ 0 \\ 1 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 0 \\ 0 \\ 0 \\ 1 \\ 1 \\ 0 \end{pmatrix}$$



$$\begin{array}{c} x = \\ \left(\begin{array}{c} 1 \\ 0 \\ 1 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 1 \\ 0 \end{array} \right) \end{array} \quad \begin{array}{c} x_1 = \\ \left(\begin{array}{c} 1 \\ 0 \\ 1 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 1 \\ 0 \end{array} \right) \end{array} \quad \begin{array}{c} x_2 = \\ \left(\begin{array}{c} 1 \\ 0 \\ 1 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 1 \\ 0 \end{array} \right) \end{array} \quad \begin{array}{c} x_3 = \\ \left(\begin{array}{c} 1 \\ 0 \\ 1 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 1 \\ 0 \end{array} \right) \end{array} \quad \begin{array}{c} x_{r1} = \\ \left(\begin{array}{c} 1 \\ 0 \\ 1 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 1 \\ 0 \end{array} \right) \end{array}$$



$$\begin{aligned}x &= \begin{pmatrix} 1 \\ 0 \\ 1 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 0 \\ 0 \\ 0 \\ 1 \\ 1 \\ 0 \end{pmatrix} & x_1 &= \begin{pmatrix} 1 \\ 0 \\ 1 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 0 \\ 0 \\ 0 \\ 1 \\ 1 \\ 0 \end{pmatrix} & x_2 &= \begin{pmatrix} 1 \\ 0 \\ 1 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 0 \\ 0 \\ 0 \\ 1 \\ 1 \\ 0 \end{pmatrix} & x_3 &= \begin{pmatrix} 1 \\ 0 \\ 1 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 0 \\ 0 \\ 0 \\ 1 \\ 1 \\ 0 \end{pmatrix} & x_{r1} &= \begin{pmatrix} 1 \\ 0 \\ 1 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 0 \\ 0 \\ 0 \\ 1 \\ 1 \\ 0 \end{pmatrix} & x_{r2} &= \begin{pmatrix} 1 \\ 0 \\ 1 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 0 \\ 0 \\ 0 \\ 1 \\ 1 \\ 0 \end{pmatrix}\end{aligned}$$



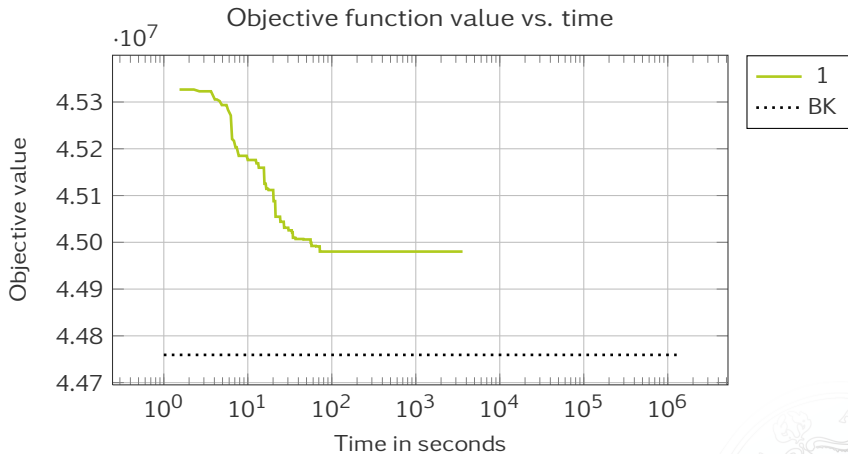
$$T_{rstu} = \max_{v,w \in \{-1,0,1\}} \frac{1}{(r-t+nv)^2 + (s-u+nw)^2}$$

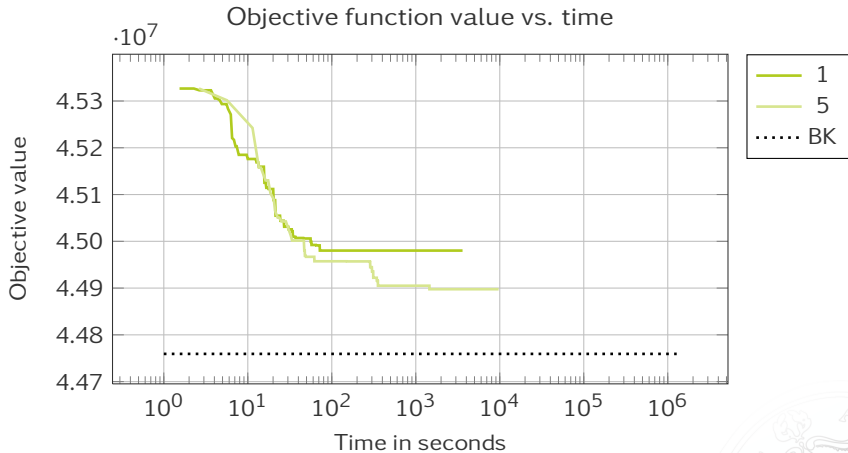
$$f_{ij} = \begin{cases} 1 & \text{if } i \leq m \text{ and } j \leq m \\ 0 & \text{otherwise} \end{cases}$$

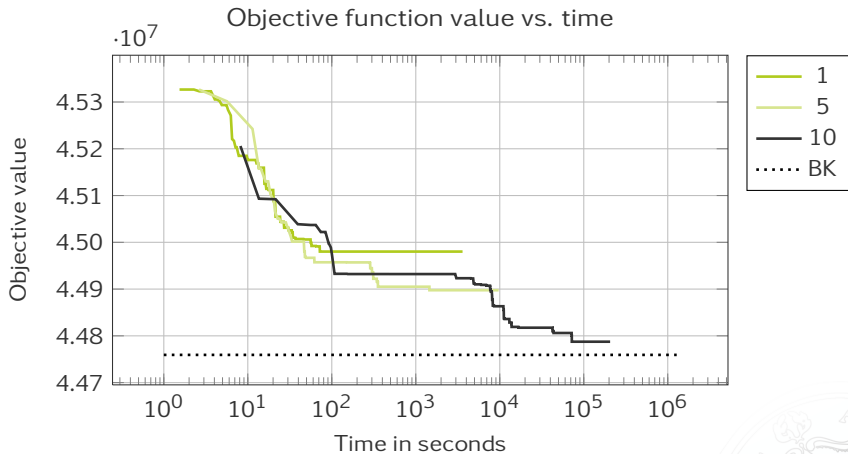
$$d_{ij} = d_{n(r-1)+s, n(t-1)+u} = T_{rstu}$$

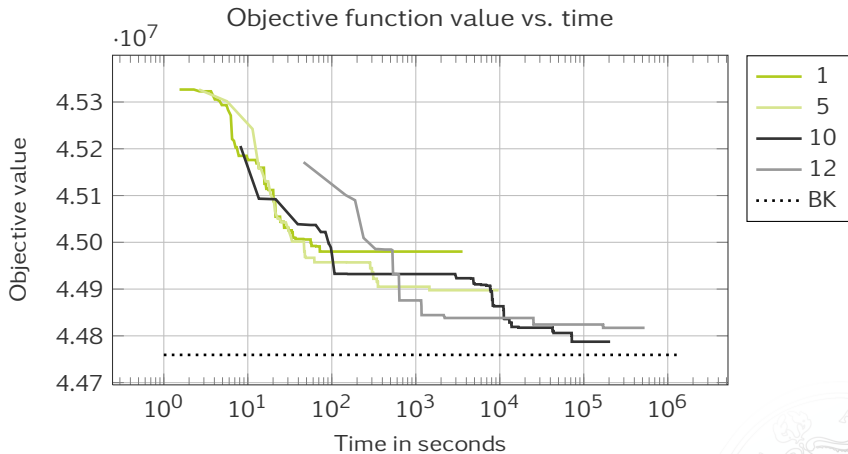
where (r, s) are the coordinates for i and (t, u) are the coordinates for j

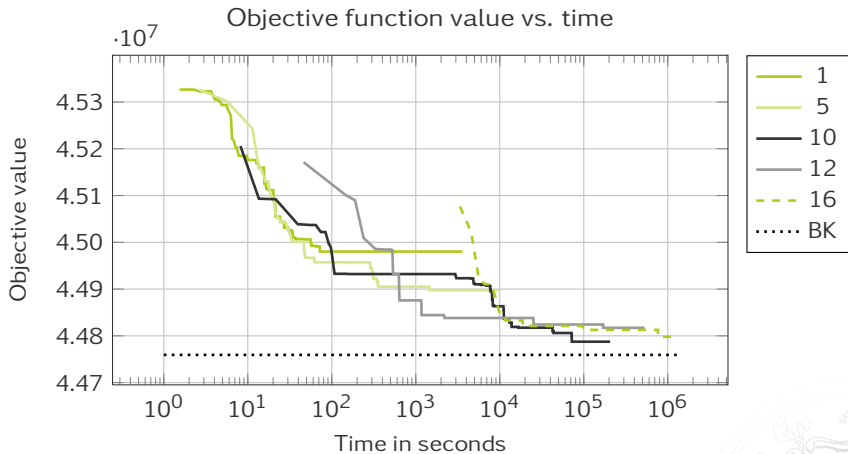


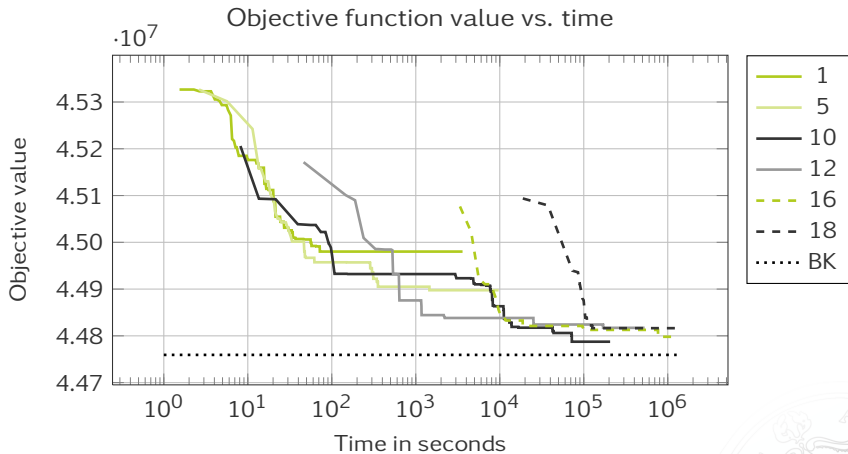


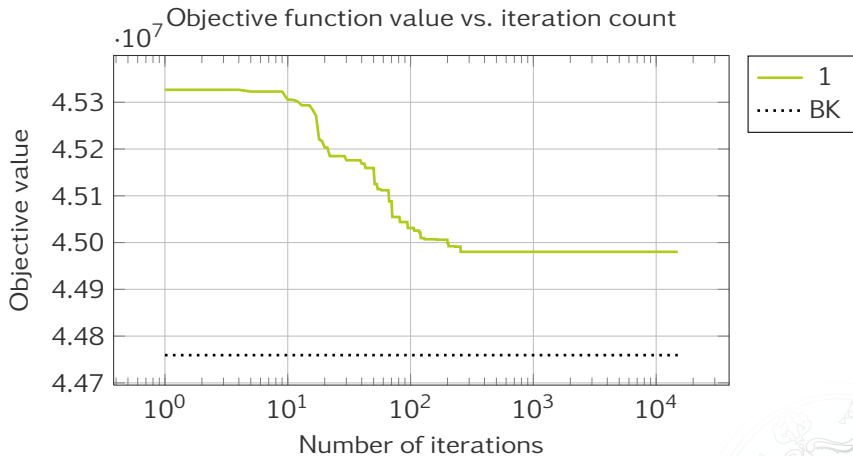


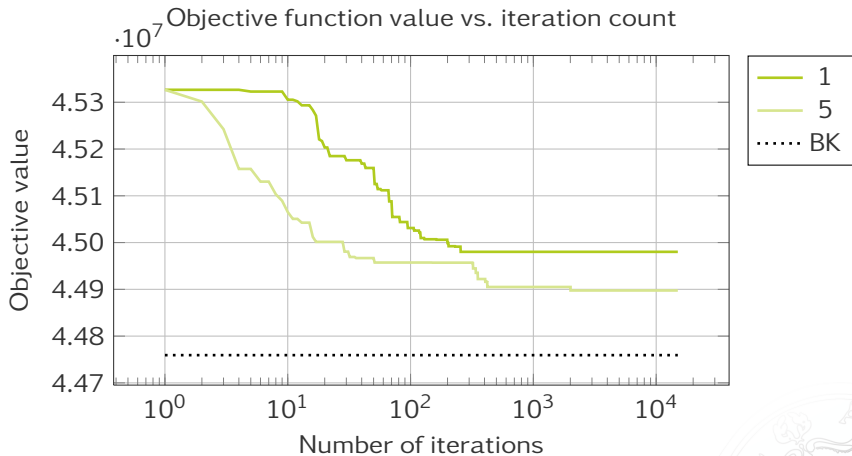


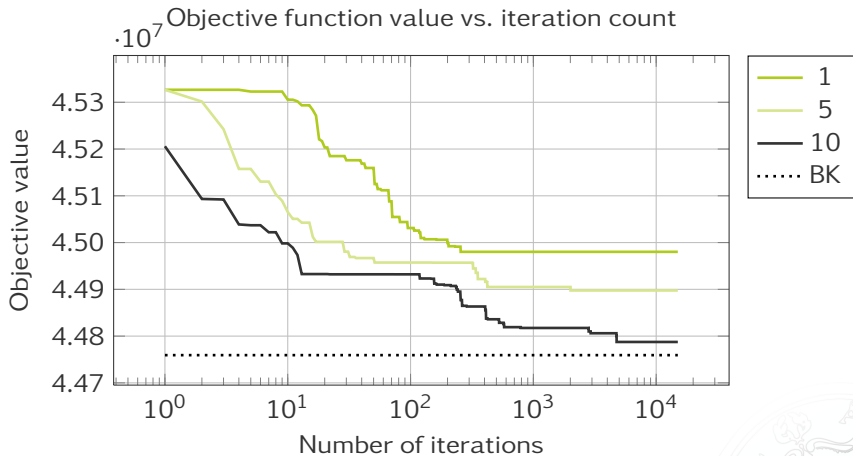


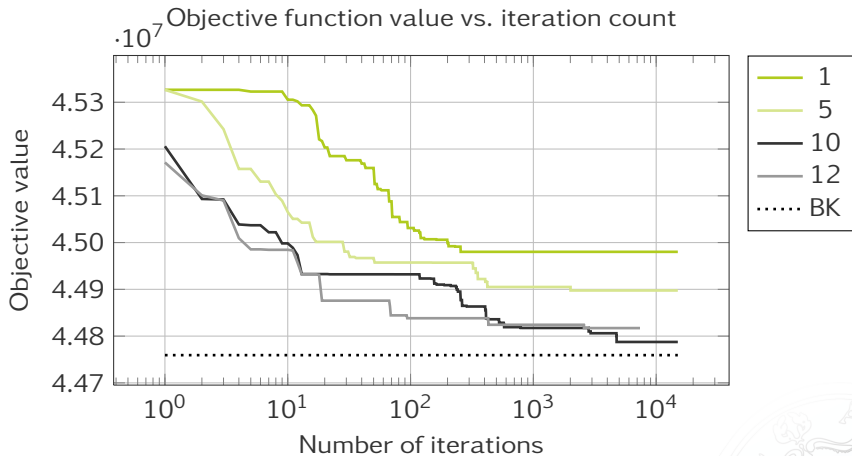


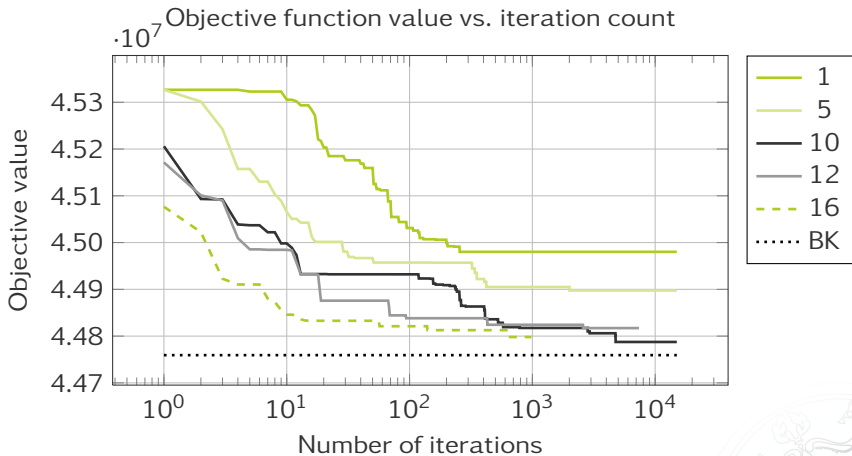


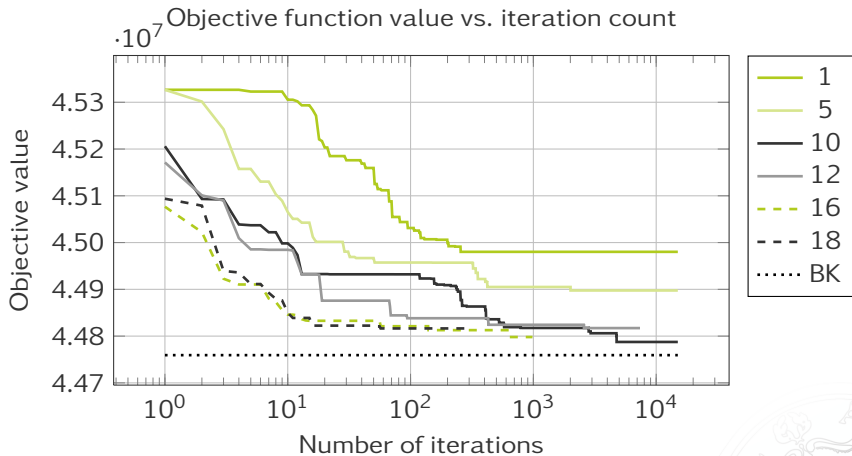


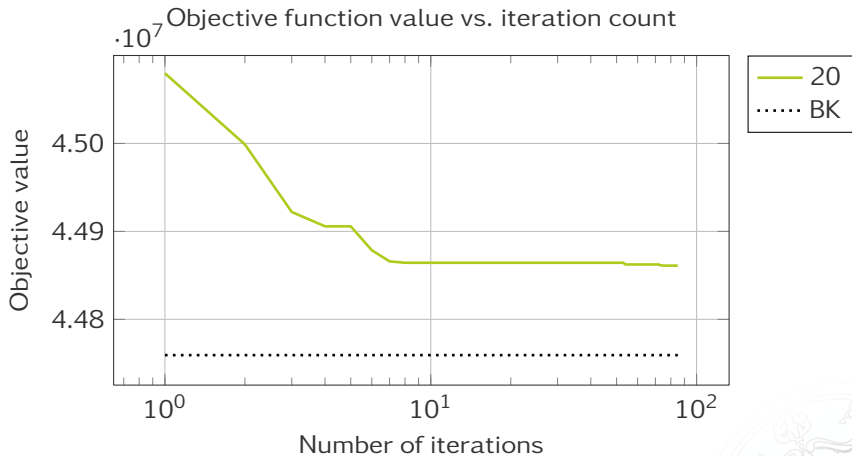


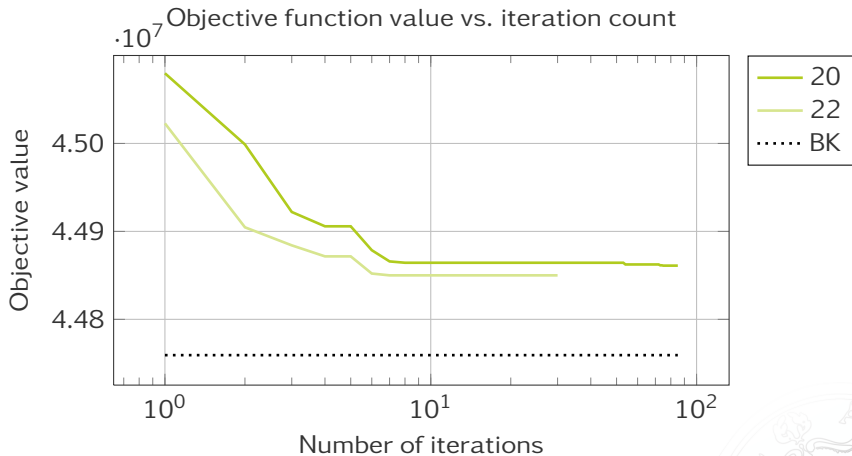


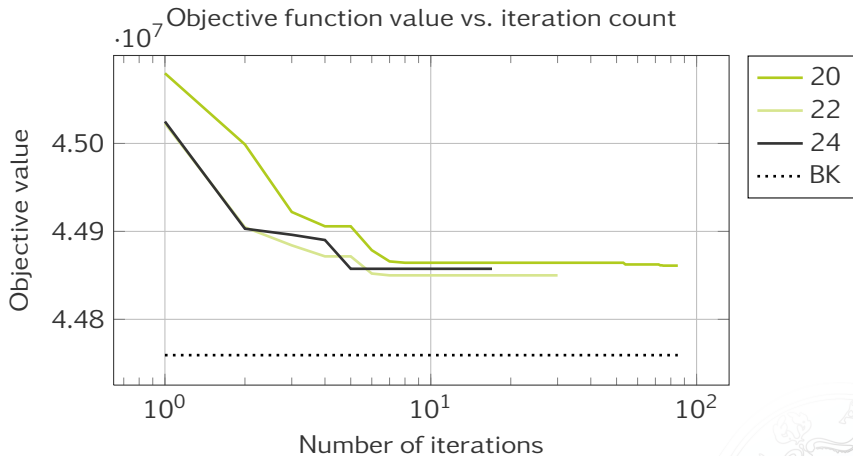


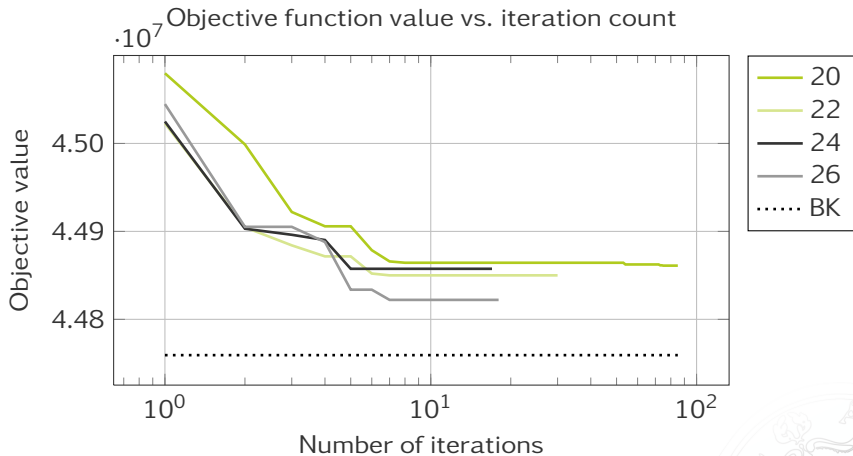


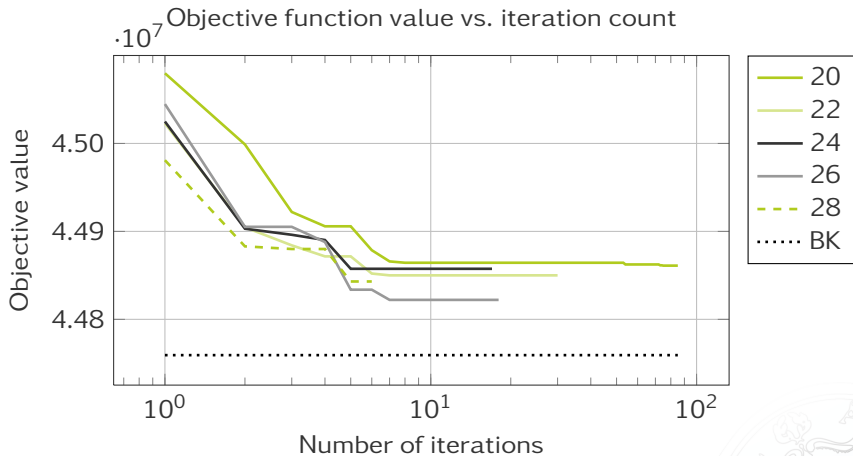


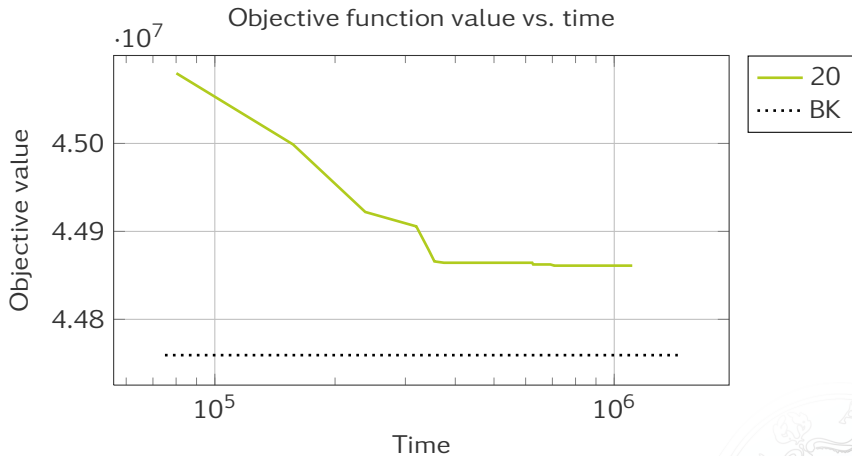


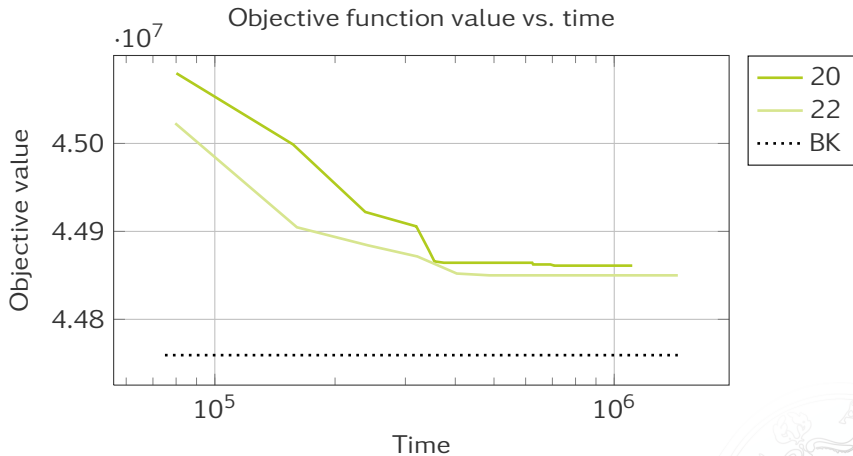


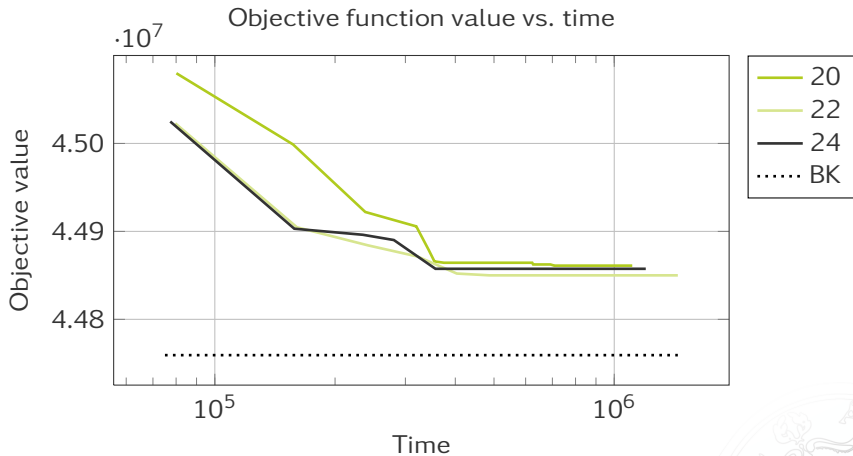


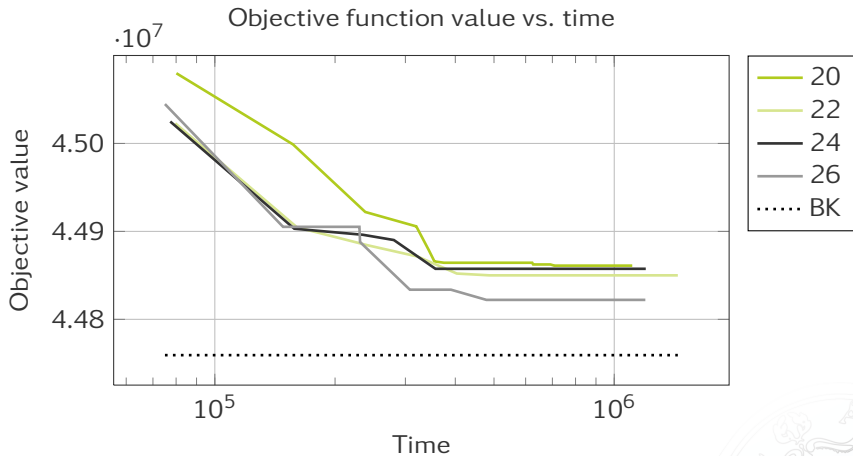


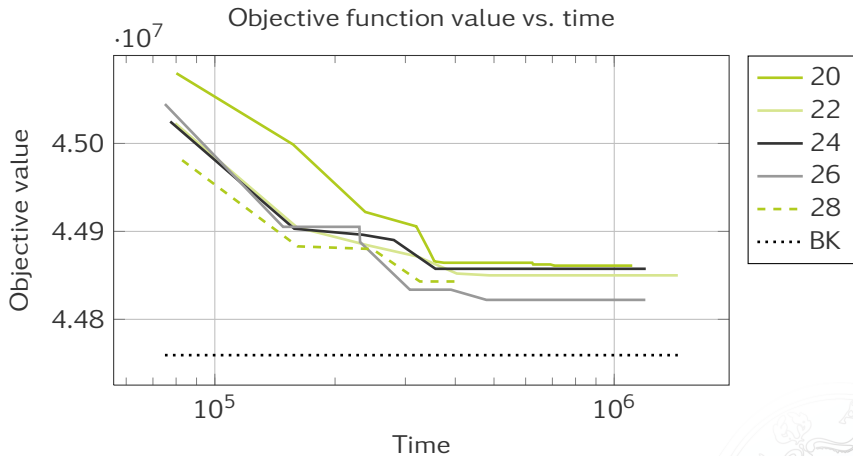












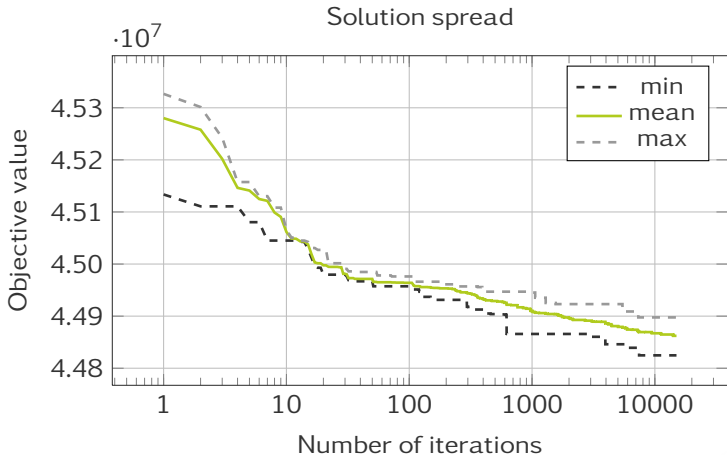


Figure 1: Spread of the solutions with $m = 5$



Some references



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The end of the presentation

Thank you for listening!

Questions?

